# Allocation problems 2A

1 
$$\begin{pmatrix} 34 & 35 & 31 \\ 26 & 31 & 27 \\ 30 & 37 & 32 \end{pmatrix} \rightarrow \text{reducing rows} \rightarrow \begin{pmatrix} 3 & 4 & 0 \\ 0 & 5 & 1 \\ 0 & 7 & 2 \end{pmatrix} \rightarrow \text{reducing columns} \rightarrow \begin{pmatrix} \frac{1}{2} & 0 & 0 \\ 0 & 3 & 2 \end{pmatrix}$$

$$\text{reducing columns} \rightarrow \begin{pmatrix} \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 2 & 1 \end{pmatrix}$$
Minimum uncovered is 1 
$$\begin{pmatrix} 4 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 2 & 1 \end{pmatrix}$$

Two solutions:

$$A-Y(35)$$
  $A-Z(31)$   
 $B-Z(27)$  or  $B-Y(31)$  cost £92  
 $C-X(30)$   $C-X(30)$ 

$$\begin{pmatrix}
34 & 37 & 32 & 32 \\
35 & 32 & 34 & 37 \\
42 & 35 & 37 & 36 \\
38 & 34 & 35 & 39
\end{pmatrix}
\rightarrow \text{reducing rows} \rightarrow
\begin{pmatrix}
2 & 5 & 0 & 0 \\
3 & 0 & 2 & 5 \\
7 & 0 & 2 & 1 \\
4 & 0 & 1 & 5
\end{pmatrix}
\rightarrow$$

→ reducing columns 
$$\begin{vmatrix}
0 & -\frac{1}{5} & 0 & 0 & 0 \\
1 & 0 & 2 & 5 \\
5 & 0 & 2 & 1 \\
2 & 0 & 1 & 5
\end{vmatrix}$$
Minimum uncovered is 1 
$$\begin{vmatrix}
0 & 6 & 0 & 0 \\
0 & 0 & 1 & 4 \\
4 & 0 & 1 & 0
\end{vmatrix}$$

Three solutions

$$\begin{pmatrix}
20 & 22 & 14 & 24 \\
20 & 19 & 12 & 20 \\
13 & 10 & 18 & 16 \\
22 & 23 & 9 & 28
\end{pmatrix}
\rightarrow \text{reducing rows} \rightarrow
\begin{pmatrix}
6 & 8 & 0 & 10 \\
8 & 7 & 0 & 8 \\
3 & 0 & 8 & 6 \\
13 & 14 & 0 & 19
\end{pmatrix}$$

reducing columns 
$$\begin{pmatrix} 3 & 8 & 0 & 4 \\ 5 & 7 & 0 & 2 \\ 0 & 0 & 8 & 0 \\ 10 & 14 & 0 & 13 \end{pmatrix}$$

Minimum uncovered element is 
$$1 \begin{pmatrix} 0 & 5 & 0 & 1 \\ 3 & 5 & 1 & 0 \\ 0 & 0 & 11 & 0 \\ 7 & 11 & 0 & 10 \end{pmatrix}$$

Minimum uncovered or element is  $1 \begin{pmatrix} 0 & 5 & 0 & 2 \\ 2 & 4 & 0 & 0 \\ 0 & 0 & 11 & 0 \\ 7 & 11 & 0 & 10 \end{pmatrix}$ 

Solution

$$J - R (20)$$

$$K - U(20)$$

$$L-S$$
 (10) cost £59

$$M-T(9)$$

4 
$$\begin{pmatrix} 85 & 95 & 97 & 87 & 80 \\ 110 & 115 & 95 & 105 & 100 \\ 90 & 95 & 86 & 93 & 105 \\ 85 & 83 & 84 & 85 & 87 \\ 100 & 100 & 105 & 120 & 95 \end{pmatrix}$$
 reducing rows  $\rightarrow$   $\begin{pmatrix} 5 & 15 & 17 & 7 & 0 \\ 15 & 20 & 0 & 10 & 5 \\ 4 & 9 & 0 & 7 & 19 \\ 2 & 0 & 1 & 2 & 4 \\ 5 & 5 & 10 & 25 & 0 \end{pmatrix}$  reducing columns  $\begin{pmatrix} 3 & 15 & 17 & 5 & 0 \\ 13 & 20 & 0 & 8 & 5 \\ 2 & 9 & 0 & 5 & 19 \\ 0 & 0 & 0 & 3 & 0 \\ 1 & 18 & 0 & 6 & 5 \\ 0 & 7 & 0 & 3 & 19 \\ 0 & 0 & 3 & 0 & 5 \\ 1 & 3 & 10 & 21 & 0 \end{pmatrix}$ 

reducing columns 
$$\begin{pmatrix}
3 & 15 & 17 & 5 & 0 \\
13 & 20 & 0 & 8 & 5 \\
2 & 9 & 0 & 5 & 19 \\
0 & 0 & 0 & 0 & 4 \\
3 & 5 & 10 & 23 & 0
\end{pmatrix}$$

### 4 (continued)

There are two solutions

$$\begin{array}{lll} D-Z & (80) & D-Y & (87) \\ E-X & (95) & E-X & (95) \\ F-V & (90) & F-V & (90) \\ G-Y & (85) & G-W & (83) & cost £450 \\ H-W & (100) & H-Z & (95) \end{array}$$

#### 5 a The initial cost matrix is shown below:

	100 m	Hurdles	200 m	400 m
Ahmed	14	21	37	64
Ben	13	22	40	68
Chang	12	20	38	70
Davina	13	21	39	74

As the table above is an n by n matrix (with n = 4), we do not need to add any additional dummy rows or columns.

The smallest numbers in rows 1, 2, 3 and 4 are 14, 13, 12 and 13. We reduce rows first by subtracting these numbers from each element in the row. The table becomes:

	100 m	Hurdles	200 m	400 m
Ahmed	0	7	23	50
Ben	0	9	27	55
Chang	0	8	26	58
Davina	0	8	26	61

The smallest numbers in columns 1, 2, 3 and 4 are 0, 7, 23 and 50. We reduce columns by subtracting these numbers from each element in the column. Therefore, the reduced cost matrix is:

	100 m	Hurdles	200 m	400 m
Ahmed	0	0	0	0
Ben	0	2	4	5
Chang	0	1	3	8
Davina	0	1	3	11

**5 b** We now apply the Hungarian algorithm.

We can cover all the zeros in two lines as follows, so the solution is not optimal.

	100 m	Hurdles	200 m	400 m
-Ahmed	· •	0	<del>0</del>	
Ben	0	2	4	5
Chang	ø	1	3	8
Davina	0	1	3	11

	100 m	Hurdles	200 m	400 m
Ahmed	0		0	0
Ben	0	2	4	5
Chang	0	1	3	8
Davina	0	1	3	11

The smallest uncovered element is 1, so we augment the matrix as follows:

- Add 1 to the elements covered by two lines
- Leave the elements covered by just one line unchanged
- Subtract 1 from the uncovered elements

This gives the following matrix:

	100 m	Hurdles	200 m	400 m
Ahmed	1	0	0	0
Ben	0	1	3	4
Chang	0	0	2	7
Davina	0	0	2	10

We can cover all the zeros in three lines, so the solution is not optimal:

	100	m	Hurdles	200 m	400 m
-Ahmed	1		<del>0</del>	0	0
Ben	0		1	3	4
Chang	0		0	2	7
Davina	0		0	2	10

	100 m	Hurdles	200 m	400 m
Ahmed		0	0	0
Ben	0	1	3	4
Chang	0	0	2	7
Davina	0	0	2	10

The smallest uncovered element is 2, so we augment the matrix as follows:

- Add 2 to the elements covered by two lines
- Leave the elements covered by just one line unchanged
- Subtract 2 from the uncovered elements

This gives the following matrix:

### 5 b (continued)

	100 m	Hurdles	200 m	400 m
Ahmed	3	2	0	0
Ben	0	1	1	2
Chang	0	0	0	5
Davina	0	0	0	8

We can cover all the zeros with four lines, therefore the solution is optimal. There are two optimal solutions of:

Ahmed - 400 m (64)

Ben - 100 m (13)

Chang - Hurdles (20)

Davina - 200 m (39)

or

Ahmed - 400 m (64)

Ben - 100 m (13)

Chang - 200 m (38)

Davina - Hurdles (21)

Both solutions give a minimum time of 136 seconds.

Minimum time = 64 + 13 + 20 + 39 = 136 seconds, or

Minimum time = 64 + 13 + 38 + 21 = 136 seconds.

**c** For an *n* by *n* matrix with n = 4, exactly 4 lines are needed to cover all the zeros.

#### **6** a The initial cost matrix is shown below:

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
A	153	87	62	144	76	68
В	162	105	87	152	88	72
C	159	84	75	165	79	77
D	145	98	63	170	85	81
E	149	94	70	138	82	89
F	160	92	82	147	80	85

As the table above is an n by n matrix (with n = 6), we do not need to add any additional dummy rows or columns.

The smallest numbers in rows 1, 2, 3, 4, 5 and 6 are 62, 72, 75, 63, 70 and 80. We reduce rows first by subtracting these numbers from each element in the row. The table becomes:

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
A	91	25	0	82	14	6
В	90	33	15	80	16	0
C	84	9	0	90	4	2
D	82	35	0	107	22	18
E	79	24	0	68	12	19
F	80	12	2	67	0	5

The smallest numbers in columns 1, 2, 3, 4, 5 and 6 are 79, 9, 0, 67, 0 and 0. We reduce columns by subtracting these numbers from each element in the column. Therefore, the reduced cost matrix is:

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
A	12	16	0	15	14	6
В	11	24	15	13	16	0
C	5	0	0	23	4	2
D	3	26	0	40	22	18
E	0	15	0	1	12	19
F	1	3	2	0	0	5

We now apply the Hungarian algorithm.

We can cover all the zeros in five lines as follows, so the solution is not optimal.

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
$\boldsymbol{A}$	12	16	0	15	14	6
В	11	24	15	13	16	0
······································	5	0		23	4	2
D	3	26	0	40	22	18
-E	0	15	0	1	12	19
- <b>F</b>	1	3	2		Ө	5

## 6 a (continued)

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
$\boldsymbol{A}$	12	16	0	15	14	6
В	11	24	15	13	16	0
C	5	0	0	23	4	2
D	3	26	0	40	22	18
E	0	15	0	1	12	<u>(9)</u>
F	1	3	2	0	0	5

The smallest uncovered element is 3, so we augment the matrix as follows:

- Add 3 to the elements covered by two lines circled in red.
- Leave the elements covered by just one line unchanged
- Subtract 3 from the uncovered elements

This gives the following matrix:

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
A	9	13	0	12	11	6
В	8	21	15	10	13	0
C	5	0	3	23	4	5
D	0	23	0	37	19	18
E	0	15	3	1	12	22
F	1	3	5	0	0	8

We can cover all the zeros in five lines as follows, so the solution is not optimal.

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
$\boldsymbol{A}$	9	13	0	12	11	6
B		21	1-5	10	13	0
C	<del>5</del>	0	3	23	4	5
D	0	23	0	37	19	18
E	0	15	3	1	12	22
F		3	5	0	0	

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
$\boldsymbol{A}$	9	13	0	12	11	6
В	8	21	13	10	13	0
C	<u>(5)</u>	0	3	23	4	5
D	0	23	0	37	19	18
E	0	15	3	1	12	22
F	1	3	5	0	0	8

### 6 a (continued)

The smallest uncovered element is 1, so we augment the matrix as follows:

- Add 1 to the elements covered by two lines circled in red
- Leave the elements covered by just one line unchanged
- Subtract 1 from the uncovered elements

This gives the following matrix:

	Beech	Elm	Eucalyptus	Oak	Olive	Willow
A	9	12	0	11	10	5
В	9	21	16	10	13	0
C	6	0	4	23	4	5
D	0	22	0	36	18	17
E	0	14	3	0	11	21
F	2	3	6	0	0	8

We can cover all the zeros with six lines, so we have found our optimal solution of:

A – Eucalyptus (£62)

B – Willow (£72)

C – Elm (£84)

D – Beech (£145)

E – Oak (£138)

F – Olive (£80)

**b** Minimum cost = 62 + 72 + 84 + 145 + 138 + 80 = £581.